

Valid Argument Forms

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An **argument** is a sequence of statements: one or more **premises**, followed by a **conclusion**. An argument is **valid** if whenever all the premises are true, the conclusion is also true — equivalently, if

$$(\text{Premise}_1 \wedge \text{Premise}_2 \wedge \cdots \wedge \text{Premise}_n) \rightarrow \text{Conclusion}$$

is a tautology.

Validity is about the argument’s *form*, not the truth of its statements: an argument can be valid with a false premise or a false conclusion. It can also have true premises and a true conclusion and still be invalid, if the truth of the conclusion doesn’t actually follow from the premises.

Provided in the exam. The named forms below are printed on the *MATH1061/MATH7861 Examination Formula Page*, alongside the logical-equivalence-laws¹ — so recognising which form an argument matches, and naming it, is the assessed skill rather than recalling the list.

Why is an argument automatically valid if it has a false premise? If any premise is false, the conjunction of the premises is false — and a false statement implies anything vacuously ($\mathbf{c} \rightarrow X$ is always T). So $(\text{Premise}_1 \wedge \cdots \wedge \text{Premise}_n) \rightarrow \text{Conclusion}$ is trivially a tautology regardless of the conclusion.

¹[courses/math1061/logical-equivalence-laws.pdf](#)

Determining validity

- **Truth table method:** build a truth table for the premises and the conclusion. The argument is valid iff every row where all premises are T also has the conclusion T.
- **Rules of inference method:** chain the laws of logical equivalence (see logical-equivalence-laws²) and the valid argument forms below to derive the conclusion from the premises.
- **Invalidity method:** assume all the premises are true and the conclusion is false, then reason out what that assumption forces each variable to be. Reaching a contradiction (some variable forced to be both T and F) means the assumption was impossible — the argument is **valid**. Reaching a consistent assignment instead gives a counterexample — the argument is **invalid**. (This directed search is usually faster than building the full truth table once there are 3+ variables.)

Valid argument forms

Name	Form(s)
Modus Ponens	$p \rightarrow q, p \therefore q$
Modus Tollens	$p \rightarrow q, \sim q \therefore \sim p$
Generalization	$p \therefore p \vee q$; or $q \therefore p \vee q$
Specialization	$p \wedge q \therefore p$; or $p \wedge q \therefore q$
Conjunction	$p, q \therefore p \wedge q$
Elimination	$p \vee q, \sim q \therefore p$; or $p \vee q, \sim p \therefore q$
Transitivity	$p \rightarrow q, q \rightarrow r \therefore p \rightarrow r$
Proof by division into cases	$p \vee q, p \rightarrow r, q \rightarrow r \therefore r$
Contradiction rule	$\sim p \rightarrow \mathbf{c} \therefore p$

These don't need to be memorised — they're a reference to check your reasoning against, not a checklist to recite.

See also

- determining-argument-validity³ — worked example applying both methods above
- logical-equivalence-laws⁴
- conditional-statements⁵
- practice-problems⁶ — §4 Valid and Invalid Arguments, with full worked solutions

²courses/math1061/logical-equivalence-laws.pdf

³courses/math1061/determining-argument-validity.pdf

⁴courses/math1061/logical-equivalence-laws.pdf

⁵courses/math1061/conditional-statements.pdf

⁶courses/math1061/practice-problems.pdf