

Python Recursion

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A **recursive function** is a function that calls itself. Every recursion needs at least one of each:

- **Base case:** a case defined outright, with no further recursive call (`if base_case: return constant`).
- **Recursive step:** a line where the function calls itself on “smaller” input, making progress toward a base case.

Canonical example: factorial

```
def fact(n: int) -> int:
    if not n:                # base case
        return 1
    return n * fact(n-1)     # recursive step
```

Evaluation has two phases:

- **Winding:** recursive calls build up a stack of pending computations ($4 \times (3 \times (2 \times (1 \times \text{fact}(0))))$).
- **Unwinding:** once the base case is hit, each pending computation resolves from the inside out.

If a base case is never reached, the recursion “bottoms out” with a `RecursionError: maximum recursion depth exceeded` once Python’s call stack limit (1000 by default) is hit. This limit can be raised with `sys.setrecursionlimit(n)`, which is sometimes necessary for algorithms that legitimately need deeper recursion (as opposed to a genuine infinite recursion bug).

Choosing base cases

Type	Typical base case
<code>int</code> (counting down)	0
<code>list</code>	<code>[]</code>
<code>str</code>	<code>''</code> (empty string)

For less obvious base cases, consider the smallest example that still needs one recursive call (the *singleton* case), and work out what value the base case must return for that example to behave correctly.

Linear vs. tree recursion

Most recursive functions make a single recursive call per invocation (linear recursion, e.g. `fact`, `is_palindrome`). Some make more than one (tree recursion), e.g. checking whether any sublist sums to a target:

```
def sublist_sum(xs: list[int], target: int) -> bool:
    if not xs:
        return not target
    return sublist_sum(xs[1:], target-xs[0]) or sublist_sum(xs[1:], target)
```

Tree recursion can recompute the same subproblem many times. **Dynamic programming** — caching previous results — avoids this when subproblems overlap heavily (the classic example being naive recursive Fibonacci, which is exponential without a cache but linear with one).

Relationship to induction

Recursion is essentially the Principle of Mathematical Induction realised as code: a base case (the $P(0)$ case) plus a step that reduces a general case to a smaller one already known to work (the $P(n) \implies P(\text{succ}(n))$ step).

See [python-scope¹](#) for how local variables/recursive calls interact with scope, and [python-composition²/python-inheritance³](#) for other structuring tools. See [towers-of-hanoi⁴](#) for a worked induction/recursion example.

¹[courses/csse1001/python-scope.pdf](#)

²[courses/csse1001/python-composition.pdf](#)

³[courses/csse1001/python-inheritance.pdf](#)

⁴[courses/csse1001/towers-of-hanoi.pdf](#)